

# Planning and Optimization

## B3. General Regression, Part I

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## B3.1 Regressing State Variables

## B3.2 Regressing Formulas Through Effects

## B3.3 Summary

## Regression for General Planning Tasks

- ▶ With disjunctions and conditional effects, things become more tricky. How to regress  $a \vee (b \wedge c)$  with respect to  $\langle q, d \triangleright b \rangle$ ?
- ▶ In this chapter, we show how to regress **general sets of states** through **general operators**.
- ▶ We extensively use the idea of representing sets of states as formulas.

## B3.1 Regressing State Variables

## Regressing State Variables: Motivation

Key question for general regression:

- ▶ Assume we are applying an operator with effect  $e$ .
- ▶ What must be true in the **predecessor state** for state variable  $v$  to be true in the **successor state**?

If we can answer this question, a general definition of regression is only a small additional step.

## Regressing State Variables: Key Idea

Assume we are in state  $s$  and apply effect  $e$  to obtain successor state  $s'$ .

State variable  $v$  is true in  $s'$  if

- ▶ effect  $e$  **makes it true**, or
- ▶ it **remains true**, i.e., it is true in  $s$  and not made false by  $e$ .

## Regressing a State Variable Through an Effect

Definition (Regressing a State Variable Through an Effect)

Let  $e$  be an effect, and let  $v$  be a state variable.

The **regression of  $v$  through  $e$** , written  $regr_e(v)$ , is defined as the following logical formula:

$$regr_e(v) = effcond_v(\tilde{e}) \vee (v \wedge \neg effcond_{\neg v}(\tilde{e})),$$

where  $\tilde{e}$  is  $e$  converted to normal form.

Questions:

- ▶ Is this well-defined?
- ▶ Why do we require normal form?

## Regressing State Variables: Example

Example

Let  $e = (b \triangleright a) \wedge (c \triangleright \neg a) \wedge b \wedge \neg d$ .

$\rightsquigarrow \tilde{e} = (b \triangleright a) \wedge ((c \wedge \neg b) \triangleright \neg a) \wedge (\top \triangleright b) \wedge (\top \triangleright \neg d)$ .

$v$	$regr_e(v)$
$a$	$b \vee (a \wedge \neg(c \wedge \neg b)) \equiv b \vee (a \wedge \neg c)$
$b$	$\top \vee (b \wedge \neg \perp) \equiv \top$
$c$	$\perp \vee (c \wedge \neg \perp) \equiv c$
$d$	$\perp \vee (d \wedge \neg \top) \equiv \perp$

Reminder:  $regr_e(v) = effcond_v(\tilde{e}) \vee (v \wedge \neg effcond_{\neg v}(\tilde{e}))$

## Regressing State Variables: Correctness (1)

### Lemma (Correctness of $\text{regr}_e(v)$ )

Let  $v$  be a state variable,  $o$  an operator and  $s$  a state in which  $o$  is applicable.

Then  $s \models \text{regr}_{\text{eff}(o)}(v)$  iff  $s[o] \models v$ .

### Proof.

Let  $\tilde{e}$  be  $\text{eff}(o)$  converted to normal form.

( $\Rightarrow$ ): We know  $s \models \text{regr}_{\text{eff}(o)}(v)$ , and hence  $s \models \text{effcond}_v(\tilde{e}) \vee (v \wedge \neg \text{effcond}_{\neg v}(\tilde{e}))$ .

Do a case analysis on the two disjuncts.

**Case 1:**  $s \models \text{effcond}_v(\tilde{e})$ . Then  $v \in [\tilde{e}]_s$  and hence  $s[o] \models v$ .

**Case 2:**  $s \models (v \wedge \neg \text{effcond}_{\neg v}(\tilde{e}))$ . Then  $s \models v$  and  $\neg v \notin [\tilde{e}]_s$ , and hence  $s[o] \models v$ . ...

## Regressing State Variables: Correctness (2)

### Proof (continued).

( $\Leftarrow$ ): Proof by contraposition.

We show that if  $\text{regr}_{\text{eff}(o)}(v)$  is **false** in  $s$ , then  $v$  is **false** in  $s[o]$ .

- ▶ By prerequisite,  $s \not\models \text{effcond}_v(\tilde{e}) \vee (v \wedge \neg \text{effcond}_{\neg v}(\tilde{e}))$ .
- ▶ Hence  $s \models \neg \text{effcond}_v(\tilde{e}) \wedge (\neg v \vee \text{effcond}_{\neg v}(\tilde{e}))$ .
- ▶ From the first conjunct, we get  $s \models \neg \text{effcond}_v(\tilde{e})$ , which implies  $s \not\models \text{effcond}_v(\tilde{e})$  and therefore  $v \notin [\tilde{e}]_s$ .
- ▶ From the second conjunct, we get  $s \models \neg v \vee \text{effcond}_{\neg v}(\tilde{e})$ .
- ▶ **Case 1:**  $s \models \neg v$ . Then  $v$  is false before applying  $o$  and remains false, so  $s[o] \not\models v$ .
- ▶ **Case 2:**  $s \models \text{effcond}_{\neg v}(\tilde{e})$ . Then  $v$  is deleted by  $o$  and not simultaneously added, so  $s[o] \not\models v$ .

□

## B3.2 Regressing Formulas Through Effects

## Regressing Formulas Through Effects: Idea

- ▶ We can now generalize regression from state variables to general formulas over state variables.
- ▶ The basic idea is to replace **every occurrence** of every state variable  $v$  by  $\text{regr}_e(v)$  as defined in the previous section.
- ▶ The following definition makes this more formal.

## Regressing Formulas Through Effects: Definition

### Definition (Regressing a Formula Through an Effect)

Let  $e$  be an effect, and let  $\varphi$  be a formula over state variables.

The **regression of  $\varphi$  through  $e$** , written  $\text{regr}_e(\varphi)$ , is defined as the following logical formula:

$$\text{regr}_e(\top) = \top$$

$$\text{regr}_e(\perp) = \perp$$

$$\text{regr}_e(v) = \text{effcond}_v(\tilde{e}) \vee (v \wedge \neg \text{effcond}_{\neg v}(\tilde{e}))$$

$$\text{regr}_e(\neg\psi) = \neg \text{regr}_e(\psi)$$

$$\text{regr}_e(\psi \vee \chi) = \text{regr}_e(\psi) \vee \text{regr}_e(\chi)$$

$$\text{regr}_e(\psi \wedge \chi) = \text{regr}_e(\psi) \wedge \text{regr}_e(\chi),$$

where  $\tilde{e}$  is  $e$  converted to normal form.

## Regressing Formulas Through Effects: Example

### Example

Let  $e = (b \triangleright a) \wedge (c \triangleright \neg a) \wedge b \wedge \neg d$ .

Recall:

$$\triangleright \text{regr}_e(a) \equiv b \vee (a \wedge \neg c)$$

$$\triangleright \text{regr}_e(b) \equiv \top$$

$$\triangleright \text{regr}_e(c) \equiv c$$

$$\triangleright \text{regr}_e(d) \equiv \perp$$

We get:

$$\begin{aligned} \text{regr}_e((a \vee d) \wedge (c \vee d)) &\equiv ((b \vee (a \wedge \neg c)) \vee \perp) \wedge (c \vee \perp) \\ &\equiv (b \vee (a \wedge \neg c)) \wedge c \\ &\equiv b \wedge c \end{aligned}$$

## Regressing Formulas Through Effects: Correctness (1)

### Lemma (Correctness of $\text{regr}_e(\varphi)$ )

Let  $\varphi$  be a logical formula,  $o$  an operator and  $s$  a state in which  $o$  is applicable.

Then  $s \models \text{regr}_{\text{eff}(o)}(\varphi)$  iff  $s[o] \models \varphi$ .

## Regressing Formulas Through Effects: Correctness (2)

### Proof.

The proof is by structural induction on  $\varphi$ .

**Induction hypothesis:**  $s \models \text{regr}_{\text{eff}(o)}(\psi)$  iff  $s[o] \models \psi$ .

**Base case  $\varphi = \top$ :**

We have  $\text{regr}_{\text{eff}(o)}(\top) = \top$ , and  $s \models \top$  iff  $s[o] \models \top$  is correct.

**Base case  $\varphi = \perp$ :**

We have  $\text{regr}_{\text{eff}(o)}(\perp) = \perp$ , and  $s \models \perp$  iff  $s[o] \models \perp$  is correct.

**Base case  $\varphi = v$ :**

We have  $s \models \text{regr}_{\text{eff}(o)}(v)$  iff  $s[o] \models v$  from the previous lemma.

...

## Regressing Formulas Through Effects: Correctness (3)

Proof (continued).

Inductive case  $\varphi = \neg\psi$ :

$$\begin{aligned} s \models \text{regr}_{\text{eff}(o)}(\neg\psi) &\text{ iff } s \models \neg \text{regr}_{\text{eff}(o)}(\psi) \\ &\text{ iff } s \not\models \text{regr}_{\text{eff}(o)}(\psi) \\ &\text{ iff } s[o] \not\models \psi \\ &\text{ iff } s[o] \models \neg\psi \end{aligned}$$

Inductive case  $\varphi = \psi \vee \chi$ :

$$\begin{aligned} s \models \text{regr}_{\text{eff}(o)}(\psi \vee \chi) &\text{ iff } s \models \text{regr}_{\text{eff}(o)}(\psi) \vee \text{regr}_{\text{eff}(o)}(\chi) \\ &\text{ iff } s \models \text{regr}_{\text{eff}(o)}(\psi) \text{ or } s \models \text{regr}_{\text{eff}(o)}(\chi) \\ &\text{ iff } s[o] \models \psi \text{ or } s[o] \models \chi \\ &\text{ iff } s[o] \models \psi \vee \chi \end{aligned}$$

Inductive case  $\varphi = \psi \wedge \chi$ :

Like previous case, replacing “ $\vee$ ” by “ $\wedge$ ”  
and replacing “or” by “and”.

□

## B3.3 Summary

## Summary

- ▶ Regressing a **state variable** through an (arbitrary) operator must consider two cases:
  - ▶ state variables **made true** (by add effects)
  - ▶ state variables **remaining true** (by absence of delete effects)
- ▶ Regression of state variables can be generalized to arbitrary formulas  $\varphi$  by replacing each occurrence of a state variable in  $\varphi$  by its regression.