# **Chapter 7** Stochastic Local Search

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# Motivation

### *n*-queens with Backtracking:



- > guarantees to find all solutions
- reaches limit for big problems:
   Best backtracking methods
   solve up to 100-queens
- > Stochastic search:
  - 1 million queens solvable
  - in less than a minute 2 29

### Systematic vs. Stochastic Search



# **Greedy Local Search**

> usually runs on complete instantiations (leaves)



- > starts in a randomly chosen instantiation
- > assignments aren't necessarily consistent

Progressing:

- > Local changes (of one variable assignment)
- Greedy, minimizing cost function (#broken constraints)

**Stopping Criterion:** 

> Assignment is consistent (const function = 0)

# **Greedy SLS: Algorithm**

#### procedure: SLS

- **Input** : A constraint network  $\mathfrak{R} = (X, D, C)$ . A cost function defined on full assignments. **Output**: A solution (no guarantee to terminate)
- **initialization:** let  $\bar{a} = (a_1, ..., a_n)$  be a random initial assignment to all variables.

#### while $\bar{a}$ is not consistent do

let  $Y = (x_i, a'_i)$  be the set of variable-value pairs that when  $x_i$  is assigned  $a'_i$ , give a maximum improvement in the cost of the assignment

pick a pair 
$$x_i, a'_i \in Y$$
.  
 $\bar{a} \leftarrow (a_1, \dots, a_{i-1}, a'_i, a_{i+1}, \dots, a_n)$  (just flip  $a_i$  to  $a'_i$ )

end

return  $\bar{a}$ 

### 4-queens with SLS:



starts in a randomly chosen instantiation

- random change of one assignment
- > minimize #broken constraints

> stop when cost function = 0



#### 4-queens with SLS:

Ŵ	3	5	2	
4	Ŵ	3	Ŵ	
3	3	6	2	
4	4	Ŵ	Ŵ	

> starts in a randomly chosen instantiation

- random change of one assignment
- > minimize #broken constraints

> stop when cost function = 0

Cost function value:

4

7 - 29

#### 4-queens with SLS:

Ŵ	Ŵ	4	2	
3	Ŵ	3	Ŵ	
2	1	5	2	
2	2	Ŵ	4	

starts in a randomly chosen instantiation

- random change of one assignment
- > minimize #broken constraints

> stop when cost function = 0

Cost function value:



#### 4-queens with SLS:

Ŵ	Ŵ	4	3	
3	2	3		
	1	2	2	
1	2	Ŵ	3	

starts in a randomly chosen instantiation

- random change of one assignment
- > minimize #broken constraints

> stop when cost function = 0

Cost function value:

1

9 - 29

# **Problem with SLS**

Search can get stuck in a *local minimum* or on a *plateau* 

2

 $\rightarrow$  Algorithm never terminates

Ŵ	1	4	2	Ŵ	1	2	2
3	Ŵ	3	Ŵ	4	2	2	Ŵ
2	1	5	2	2	Ŵ	3	3
2	2	Ŵ	4	3	2	Ŵ	3

Cost function value:

Cost function value: 1

### Plateaus & Local Minima



### 1. Plateau Search

- > Allow non-improving sideway steps
- > Problem: running in circles



### 2. Tabu search

- Store last n variable-value assignments
- > Use list to prevent backward moves



### **3. Random Restarts**

- > Restart algorithm in new random initialisation
- Can be combined with other escape-techniques
- > Suggestions for restart:
  - > when no improvement is possible
  - > after max\_flips steps without improvement (Plateau search)
  - > increase max\_flips after every improvement

> Achieve guarantee to find a solution

### 4. Constraint weighting

- > Cost function:  $F(\vec{a}) = \sum w_i * C_i(\vec{a})$
- Increasing weights of a violated constraint in local minima



## **Other improvements**

### **Problem: Undetermined Termination**

- Set a limit max\_tries for the algorithm when to stop
- but: we lose guarantee to find a solution

### **Anytime Behaviour**

- Store best assignment found so far (minimal #broken constraints)
- Return assignment when we need one (no solution)

# Random Walks

procedure: RandomWalk

- **Input** : A network  $\mathfrak{R} = (X, D, C)$ , probability p.
- **Output**: A solution iff the problem is consistent.

**start** with a random initial assignment  $\bar{a}$ .

while  $\bar{a}$  is not a solution do

- (i) **pick** a violated constraint C, randomly
- (ii) **choose** with probability p a variable-value pair  $\langle x, a' \rangle$  for  $x \in scope(C)$ , or, with probability 1 - p, choose a variable-value pair  $\langle x, a' \rangle$ that minimizes the value of the cost function when the value of x is changed to a'.
- (iii) Change x's value to a'.

end

return  $\bar{a}$ .

### Eventually hits a satisfying assignment (if exists) 17 - 29

# p and Simulated Annealing

> Optimal p values for specific problems

### Extension: Simulated Annealing

- > Decrease p over time (by "cooling the temperature")
  - > more random jumps in earlier stages
  - > more greedy progress later

## SLS + Inference

Goal: Smaller search space

- > use Inference methods as with systematic search
- > constraint propagation: performance varies
  - very helpful for removing many near-solutions
  - > not good for uniform problem structures

### Recap: Cycle-cutset decomposition



Idea: Replace systematic search on cutset with SLS

Start with random cutset assignment

Repeat:

> calculate minimal cost in trees:

$$C(z_i \rightarrow a_i) = \sum_{children \ z_j} min_{a_j \in D_{z_j}} (C(z_j \rightarrow a_j) + R(z_i \rightarrow a_i, z_j \rightarrow a_j))$$

- > assign values with minimal cost to tree variables
- > greedily optimize cutset assignment (Local Search)

### Example: Binary domains

1. Assign values to cutset variables





Set a **Root** for each tree





2. From leaves to root:



3. From root to leaves:

Assign values with minimal cost





1. Assign values to cutset variables





2. From leaves to root:



 $C(z_i \rightarrow a_i) = \sum_{\text{children } z_j} \min_{a_j \in D_{z_j}} (C(z_j \rightarrow a_j) + R(z_i \rightarrow a_i, z_j \rightarrow a_j))$  27 - 29

3. From root to leaves:

Assign values with minimal cost



## Summary

### **Stochastic Local Search**

- > Approximates systematic search
- Greedy algorithms: Techniques to escape local minima
- Random Walk: combines greedy + random choices
- Combination with Inference methods can help

- Can work very well
- > but no guarantee of termination AND finding a solution