

# Foundations of Artificial Intelligence

## 10. State-Space Search: Breadth-first Search

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# Foundations of Artificial Intelligence

March 15, 2023 — 10. State-Space Search: Breadth-first Search

## 10.1 Blind Search

## 10.2 Breadth-first Search: Introduction

## 10.3 BFS-Tree

## 10.4 BFS-Graph

## 10.5 Properties of Breadth-first Search

## 10.6 Summary

## State-Space Search: Overview

### Chapter overview: state-space search

- ▶ 5.–7. Foundations
- ▶ 8.–12. Basic Algorithms
  - ▶ 8. Data Structures for Search Algorithms
  - ▶ 9. Tree Search and Graph Search
  - ▶ 10. Breadth-first Search
  - ▶ 11. Uniform Cost Search
  - ▶ 12. Depth-first Search and Iterative Deepening
- ▶ 13.–19. Heuristic Algorithms

## 10.1 Blind Search

## Blind Search

In Chapters 10–12 we consider **blind** search algorithms:

### Blind Search Algorithms

Blind search algorithms use **no** information about state spaces apart from the black box interface.

They are also called **uninformed** search algorithms.

contrast: **heuristic** search algorithms (Chapters 13–19)

## Blind Search Algorithms: Examples

examples of blind search algorithms:

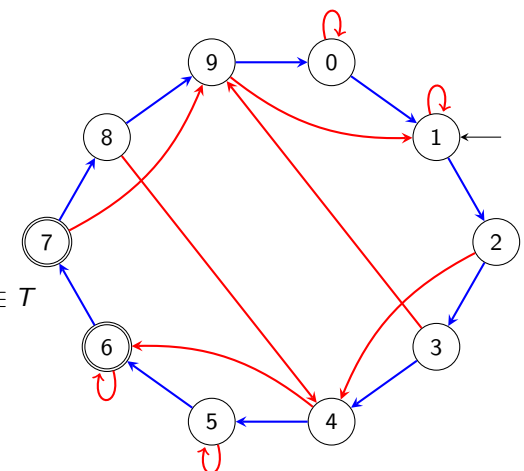
- ▶ **breadth-first search** (↪ this chapter)
- ▶ uniform cost search (↪ Chapter 11)
- ▶ depth-first search (↪ Chapter 12)
- ▶ depth-limited search (↪ Chapter 12)
- ▶ iterative deepening search (↪ Chapter 12)

## 10.2 Breadth-first Search: Introduction

## Running Example: Reminder

bounded inc-and-square:

- ▶  $S = \{0, 1, \dots, 9\}$
- ▶  $A = \{inc, sqr\}$
- ▶  $cost(inc) = cost(sqr) = 1$
- ▶  $T$  s.t. for  $i = 0, \dots, 9$ :
  - ▶  $\langle i, inc, (i + 1) \bmod 10 \rangle \in T$
  - ▶  $\langle i, sqr, i^2 \bmod 10 \rangle \in T$
- ▶  $s_l = 1$
- ▶  $S_* = \{6, 7\}$



## Idea and Example

breadth-first search:

- ▶ expands nodes **in order of generation** (FIFO)  
 ~> e.g., open list as **linked list** or **deque**
- ▶ different variants and optimizations (~> later)
  - ▶ Use a closed list?
  - ▶ When to update closed list?
  - ▶ When to perform duplicate check?
  - ▶ When to perform goal test?

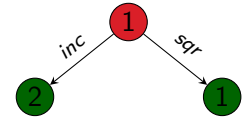


open:  $\left[ \begin{array}{c} \text{next} \\ \downarrow \\ \bullet \end{array} \right]$   
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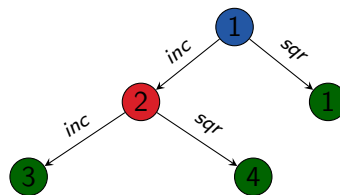


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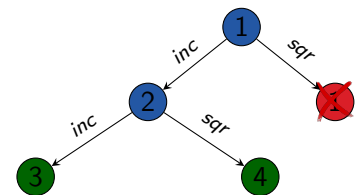


open:  $\left[ \begin{array}{c} \text{next} \\ \downarrow \\ \bullet \quad \bullet \quad \bullet \end{array} \right]$   
 closed:  $\{1, 2\}$

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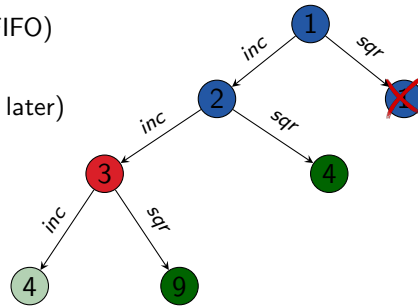


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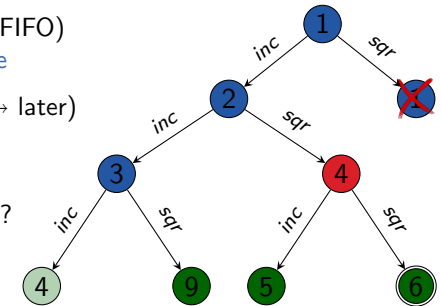
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closed:  $\{1, 2, 3\}$

## Idea and Example

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  - When to perform goal test?
- searches state space **layer by layer**
- always finds **shallowest** goal state first



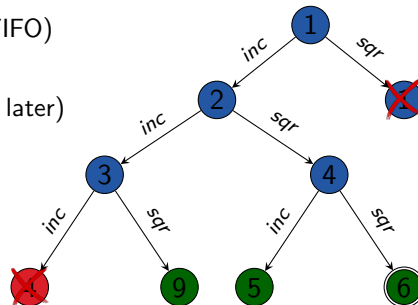
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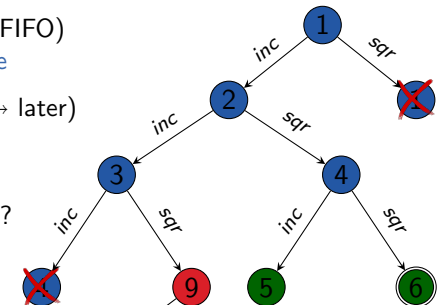
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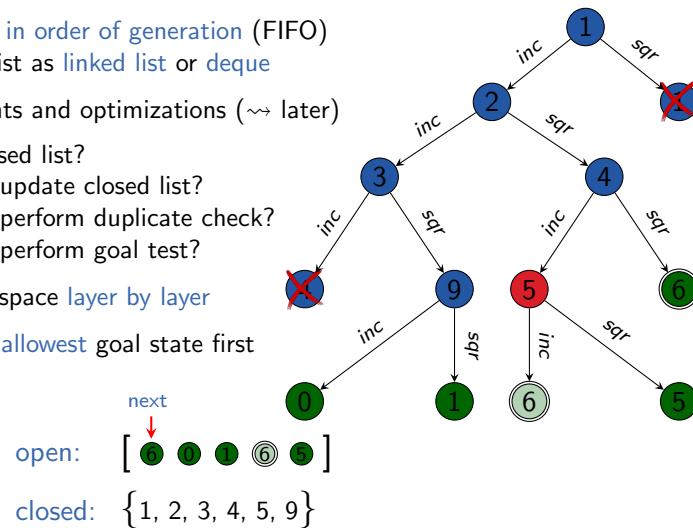
open:  $\left[ \begin{array}{c} \text{next} \\ \downarrow \\ \bullet \end{array} \right]$

closed:  $\{1, 2, 3, 4, 9\}$

## Idea and Example

breadth-first search:

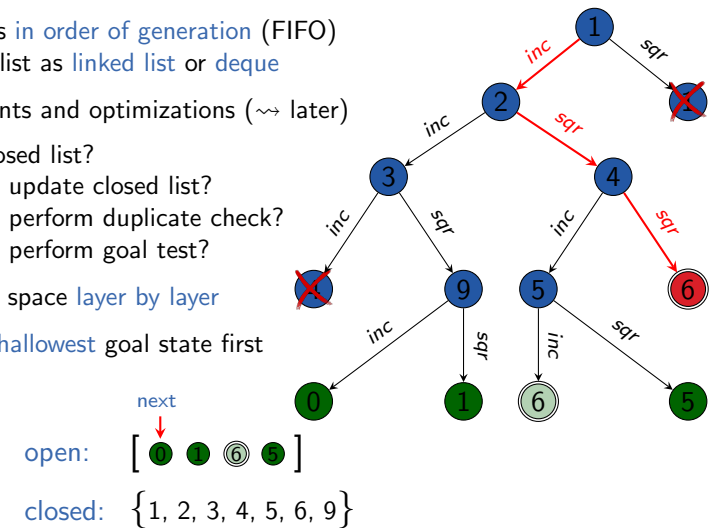
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## Breadth-first Search: Tree Search or Graph Search?

Breadth-first search can be performed

- ▶ **without duplicate elimination** (as a tree search)  
 ~> **BFS-Tree**
- ▶ **or with duplicate elimination** (as a graph search)  
 ~> **BFS-Graph**

(BFS = **breadth-first search**).

~> We consider both variants.

## 10.3 BFS-Tree

## Reminder: Generic Tree Search Algorithm

reminder from Chapter 9:

### Generic Tree Search

```

open := new OpenList
open.insert(make_root_node())
while not open.is_empty():
    n := open.pop()
    if is_goal(n.state):
        return extract_path(n)
    for each  $\langle a, s' \rangle \in \text{succ}(n.\text{state})$ :
        n' := make_node(n, a, s')
        open.insert(n')
return unsolvable

```

## BFS-Tree (1st Attempt)

breadth-first search without duplicate elimination (1st attempt):

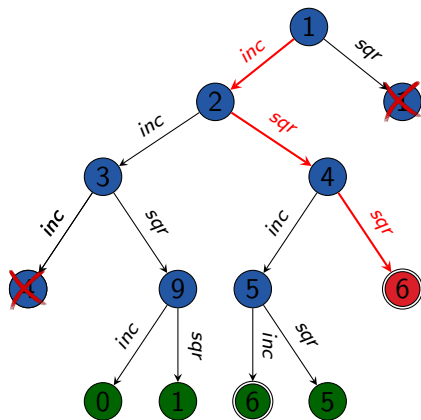
### BFS-Tree (1st Attempt)

```

open := new Deque
open.push_back(make_root_node())
while not open.is_empty():
    n := open.pop_front()
    if is_goal(n.state):
        return extract_path(n)
    for each  $\langle a, s' \rangle \in \text{succ}(n.\text{state})$ :
        n' := make_node(n, a, s')
        open.push_back(n')
return unsolvable

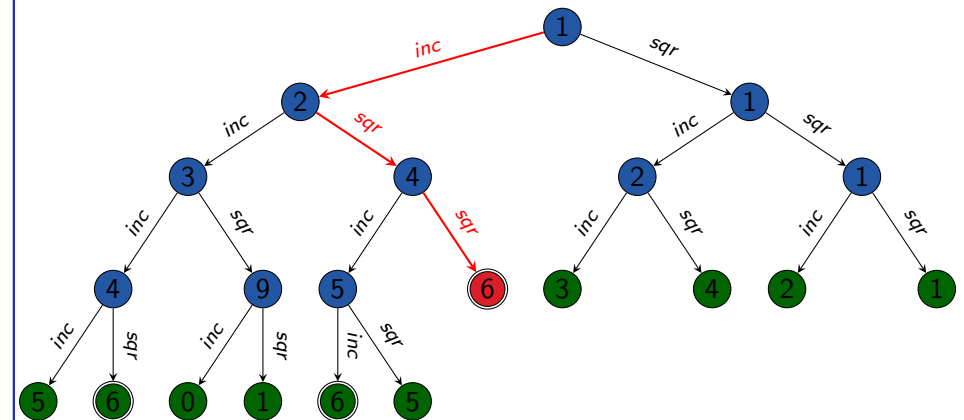
```

## Example and Discussion



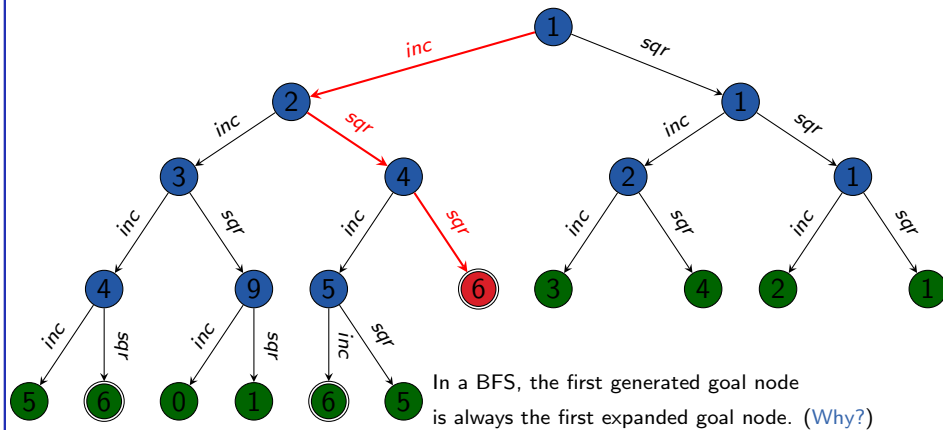
How does the search tree of the initial example differ if we run BFS-Tree (1st Attempt) instead?

## Example and Discussion

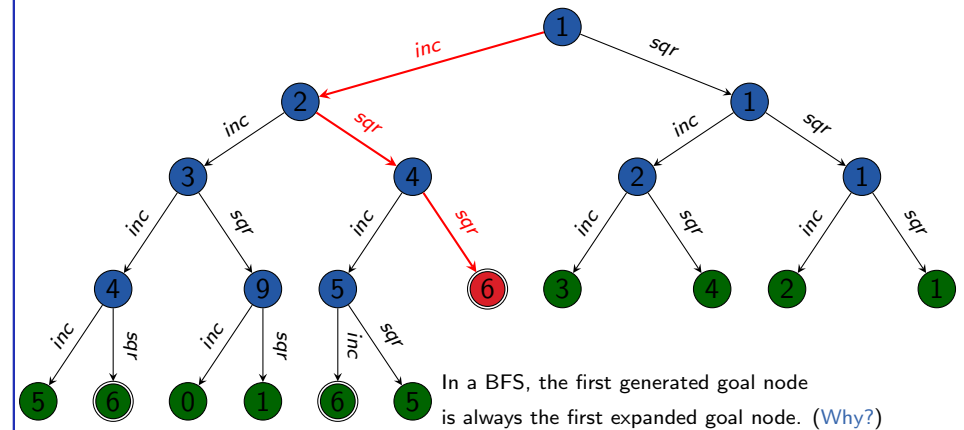


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## Example and Discussion



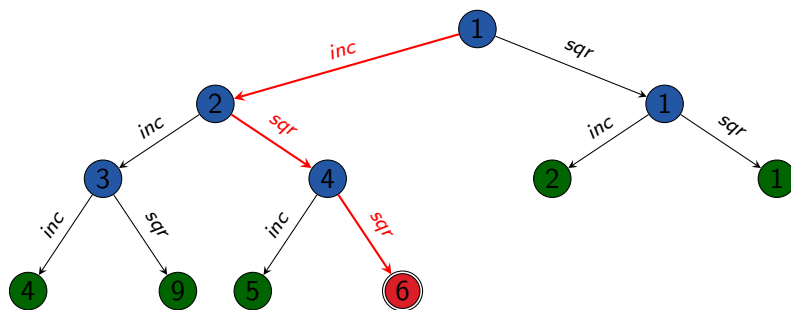
## Example and Discussion



→ It is more efficient to perform the goal test upon **generating** a node (rather than upon **expanding** it).

→ How much effort does this save?

## Example and Discussion



→ It is more efficient to perform the goal test upon **generating** a node (rather than upon **expanding** it).

→ How much effort does this save?

## BFS-Tree (2nd Attempt)

breadth-first search without duplicate elimination (2nd attempt):

## BFS-Tree (2nd Attempt)

```

open := new Queue
open.push_back(make_root_node())
while not open.is_empty():
    n := open.pop_front()
    if is_goal(n.state):
        return extract_path(n)
    for each <a, s'> in succ(n.state):
        n' := make_node(n, a, s')
        if is_goal(n'.state):
            return extract_path(n')
        open.push_back(n')
return unsolvable

```

## BFS-Tree (2nd Attempt): Discussion

Where is the bug?

## BFS-Tree (Final Version)

breadth-first search without duplicate elimination (final version):

```

BFS-Tree
if is_goal(init()):
    return ⟨⟩
open := new Deque
open.push_back(make_root_node())
while not open.is_empty():
    n := open.pop_front()
    for each ⟨a, s'⟩ ∈ succ(n.state):
        n' := make_node(n, a, s')
        if is_goal(s'):
            return extract_path(n')
        open.push_back(n')
return unsolvable

```

## 10.4 BFS-Graph

## Reminder: Generic Graph Search Algorithm

reminder from Chapter 9:

```

Generic Graph Search
open := new OpenList
open.insert(make_root_node())
closed := new ClosedList
while not open.is_empty():
    n := open.pop()
    if closed.lookup(n.state) = none:
        closed.insert(n)
        if is_goal(n.state):
            return extract_path(n)
        for each ⟨a, s'⟩ ∈ succ(n.state):
            n' := make_node(n, a, s')
            open.insert(n')
return unsolvable

```



## Adapting Generic Graph Search to Breadth-First Search

adapting the generic algorithm to breadth-first search:

- ▶ similar adaptations to BFS-Tree  
(**deque** as open list, **early goal test**)
- ▶ as closed list does not need to manage node information, a **set** data structure suffices
- ▶ for the same reasons why early goal tests are a good idea, we should perform **duplicate tests** against the closed list and **updates of the closed lists** as early as possible

## BFS-Graph (Breadth-First Search with Duplicate Elim.)

### BFS-Graph

```

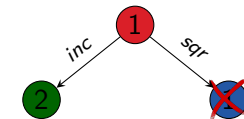
if is_goal(init()):
    return {}
open := new Deque
open.push_back(make_root_node())
closed := new HashSet
closed.insert(init())
while not open.is_empty():
    n := open.pop_front()
    for each  $\langle a, s' \rangle \in \text{succ}(n.\text{state})$ :
        n' := make_node(n, a, s')
        if is_goal(s'):
            return extract_path(n')
        if s'  $\notin$  closed:
            closed.insert(s')
            open.push_back(n')
return unsolvable
  
```

## Example



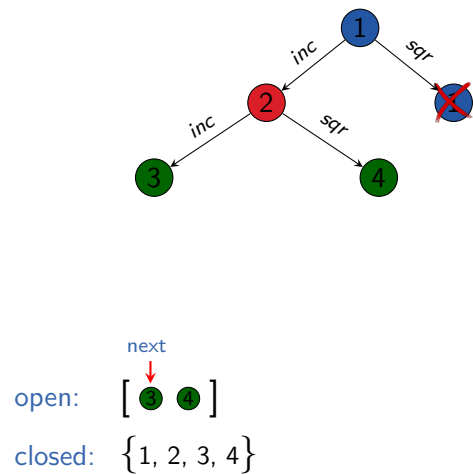
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## Example

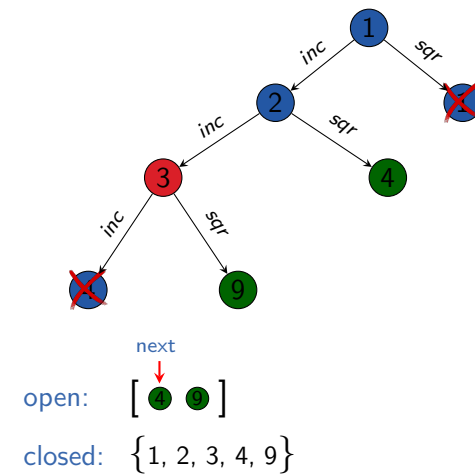


open:  $\left[ \begin{array}{c} \text{next} \\ \downarrow \\ \bullet \end{array} \right]$   
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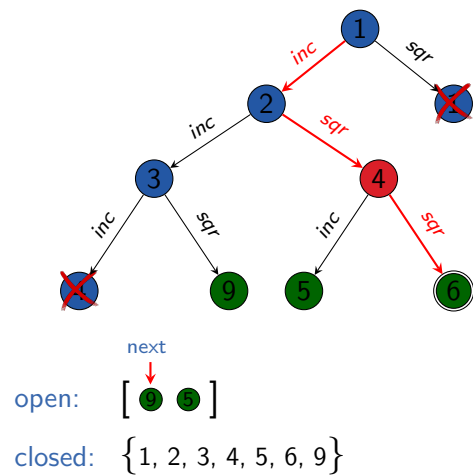
## Example



## Example



## Example



## 10.5 Properties of Breadth-first Search

## Properties of Breadth-first Search

### Properties of Breadth-first Search:

- ▶ BFS-Tree is **semi-complete**, but not **complete**. (Why?)
- ▶ BFS-Graph is **complete**. (Why?)
- ▶ BFS (both variants) is **optimal** if all actions have the same cost (Why?), but not in general (Why not?).
- ▶ complexity: **next slides**

## Breadth-first Search: Complexity

The following result applies to both BFS variants:

### Theorem (time complexity of breadth-first search)

Let  $b$  be the branching factor and  $d$  be the minimal solution length of the given state space. Let  $b \geq 2$ .

Then the **time complexity** of breadth-first search is

$$1 + b + b^2 + b^3 + \dots + b^d = O(b^d)$$

**Reminder:** we measure time complexity in generated nodes.

It follows that the **space complexity** of both BFS variants also is  $O(b^d)$  (if  $b \geq 2$ ). (Why?)

## Breadth-first Search: Example of Complexity

**example:**  $b = 13$ ; 100 000 nodes/second; 32 bytes/node

$d$	nodes	time	memory
4	30 940	0.3 s	966 KiB
6	$5.2 \cdot 10^6$	52 s	159 MiB
8	$8.8 \cdot 10^8$	147 min	26 GiB
10	$10^{11}$	17 days	4.3 TiB
12	$10^{13}$	8 years	734 TiB
14	$10^{15}$	1 352 years	121 PiB
16	$10^{17}$	$2.2 \cdot 10^5$ years	20 EiB
18	$10^{20}$	$38 \cdot 10^6$ years <b><math>38 \cdot 10^6</math> years</b>	3.3 ZiB

## Breadth-first Search: Example of Complexity

**example:**  $b = 13$ ; 100 000 nodes/second; 32 bytes/node

**Realistic numbers?**

$d$	nodes	time	memory
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## Breadth-first Search: Example of Complexity

example:  $b = 13$ ; 100 000 nodes/second; 32 bytes/node



Rubik's cube:

- ▶ branching factor:  $\approx 13$
- ▶ typical solution length: 18

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## BFS-Tree or BFS-Graph?

What is better, BFS-Tree or BFS-Graph?

advantages of BFS-Graph:

- ▶ complete
- ▶ much (!) more efficient if there are many duplicates

advantages of BFS-Tree:

- ▶ simpler
- ▶ less overhead (time/space) if there are few duplicates

### Conclusion

BFS-Graph is usually preferable, unless we know that there is a negligible number of duplicates in the given state space.

## 10.6 Summary

## Summary

- ▶ **blind search algorithm**: use no information except black box interface of state space
- ▶ **breadth-first search**: expand nodes in order of generation
  - ▶ search state space **layer by layer**
  - ▶ can be tree search or graph search
  - ▶ complexity  $O(b^d)$  with branching factor  $b$ , minimal solution length  $d$  (if  $b \geq 2$ )
  - ▶ **complete** as a graph search; **semi-complete** as a tree search
  - ▶ **optimal** with **uniform action costs**