Foundations of Artificial Intelligence
26. Constraint Satisfaction Problems: Path Consistency

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Constraint Satisfaction Problems: Overview

Chapter overview: constraint satisfaction problems:

26.1 Beyond Arc Consistency

Beyond Arc Consistency: Path Consistency

idea of arc consistency:

- ▶ For every assignment to a variable u
  there must be a suitable assignment to every other variable v.
- ▶ If not: remove values of u for which
  no suitable "partner" assignment to v exists.
- ⇝ tighter unary constraint on u

This idea can be extended to three variables (path consistency):

- ▶ For every joint assignment to variables u, v
  there must be a suitable assignment to every third variable w.
- ▶ If not: remove pairs of values of u and v for which
  no suitable "partner" assignment to w exists.
- ⇝ tighter binary constraint on u and v

German: Pfadkonsistenz

Beyond Arc Consistency: i-Consistency

general concept of i-consistency for $i \geq 2$:

- ▶ For every joint assignment to variables $v_1, \ldots, v_{i-1}$
  there must be a suitable assignment to every i-th variable v_i.
- ▶ If not: remove value tuples of $v_1, \ldots, v_{i-1}$ for which
  no suitable "partner" assignment for v_i exists.
- ⇝ tighter $(i-1)$-ary constraint on $v_1, \ldots, v_{i-1}$
- ▶ 2-consistency = arc consistency
- ▶ 3-consistency = path consistency (*)

We do not consider general i-consistency further
as larger values than $i = 3$ are rarely used
and we restrict ourselves to binary constraints in this course.

(*) usual definitions of 3-consistency vs. path consistency differ
     when ternary constraints are allowed

26.2 Path Consistency

## Path Consistency: Definition

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Definition (path consistent)

Let $\mathcal{C} = \langle V, \text{dom}, (R_{uv}) \rangle$ be a constraint network.

(a) Two different variables $u, v \in V$ are path consistent
with respect to a third variable $w \in V$ if
for all values $d_u \in \text{dom}(u), d_v \in \text{dom}(v)$ with $\langle d_u, d_v \rangle \in R_{uv}$
there is a value $d_w \in \text{dom}(w)$ with $\langle d_u, d_w \rangle \in R_{uw}$ and
$\langle d_v, d_w \rangle \in R_{vw}$.

(b) The constraint network $\mathcal{C}$ is path consistent
if for any three variables u, v, w,
the variables u and v are path consistent with respect to w.
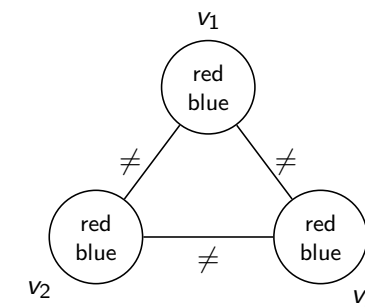

Path Consistency: Remarks

remarks:

- ▶ Even if the constraint R_{uv} is trivial, path consistency can infer nontrivial constraints between u and v.
- ▶ name "path consistency":
  path $u \to w \to v$ leads to new information on $u \to v$
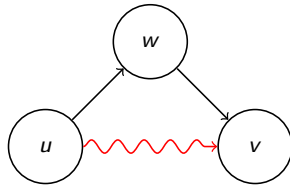
Path Consistency: Example

arc consistent, but not path consistent

Processing Variable Triples: revise-3

analogous to revise for arc consistency:

> **function** revise-3($\mathcal{C}, u, v, w$):
> $\langle V, \text{dom}, (R_{uv}) \rangle := \mathcal{C}$
> **for each** $\langle d_u, d_v \rangle \in R_{uv}$:
>   **if** there is no $d_w \in \text{dom}(w)$ with
> $\langle d_u, d_w \rangle \in R_{uw}$ **and** $\langle d_v, d_w \rangle \in R_{vw}$:
>   **remove** $\langle d_u, d_v \rangle$ from R_{uv}

input: constraint network $\mathcal{C}$ and three variables u, v, w of $\mathcal{C}$

effect: u, v path consistent with respect to w.
All violating pairs are removed from R_{uv}.

time complexity: $O(k^3)$ where k is maximal domain size

Enforcing Path Consistency: PC-2

analogous to AC-3 for arc consistency:

> **function** PC-2($\mathcal{C}$):
> $\langle V, \text{dom}, (R_{uv}) \rangle := \mathcal{C}$
> $queue := \emptyset$
> **for each** set of two variables $\{u, v\}$:
>   **for each** $w \in V \setminus \{u, v\}$:
>   insert $\langle u, v, w \rangle$ into $queue$
> **while** $queue \neq \emptyset$:
>   remove any element $\langle u, v, w \rangle$ from $queue$
>   revise-3($\mathcal{C}, u, v, w$)
>   **if** R_{uv} changed in the call to revise-3:
>   **for each** $w' \in V \setminus \{u, v\}$:
>   insert $\langle w', u, v \rangle$ into $queue$
>   insert $\langle w', v, u \rangle$ into $queue$

PC-2: Discussion

The comments for AC-3 hold analogously.

▶ PC-2 enforces path consistency

▶ proof idea: invariant of the **while** loop:
if $\langle u, v, w \rangle \notin queue$, then u, v path consistent
with respect to w

▶ time complexity $O(n^3 k^5)$ for n variables and maximal domain
size k (Why?)

26.3 Summary

Summary

▶ generalization of
arc consistency (considers pairs of variables)
to path consistency (considers triples of variables)
and i-consistency (considers i-tuples of variables)

▶ arc consistency tightens unary constraints

▶ path consistency tightens binary constraints

▶ i-consistency tightens $(i - 1)$-ary constraints

▶ higher levels of consistency more powerful
but more expensive than arc consistency