Theory of Computer Science

C1. Formal Languages and Grammars

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Theory of Computer Science

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C1.1 Introduction

Course Contents

Parts of the course:

A. background ✓
   - ▷ mathematical foundations and proof techniques

B. logic (Logik) ✓
   - ▷ How can knowledge be represented?
     How can reasoning be automated?

C. automata theory and formal languages
   (Automatentheorie und formale Sprachen)
   - ▷ What is a computation?

D. computability theory (Berechenbarkeitstheorie)
   - ▷ What can be computed at all?

E. complexity theory (Komplexitätstheorie)
   - ▷ What can be computed efficiently?

Example: Propositional Formulas

from the logic part:

Definition (Syntax of Propositional Logic)

Let A be a set of atomic propositions. The set of propositional formulas (over A) is inductively defined as follows:

- Every atom $a \in A$ is a propositional formula over A.
- If φ is a propositional formula over A, then so is its negation $\neg\varphi$.
- If φ and ψ are propositional formulas over A, then so is the conjunction $(\varphi \land \psi)$.
- If φ and ψ are propositional formulas over A, then so is the disjunction $(\varphi \lor \psi)$.

Example: Propositional Formulas

Let S_A be the set of all propositional formulas over A.

Such sets of symbol sequences (or words) are called languages.

Sought: General concepts to define such (often infinite) languages with finite descriptions.

- today: grammars
- later: automata

Example: Propositional Formulas

Example (Grammar for $S_{\{\mathrm{a},\mathrm{b},\mathrm{c}\}}$)

Grammar variables $\{\mathrm{F}, \mathrm{A}, \mathrm{N}, \mathrm{C}, \mathrm{D}\}$ with start variable F, terminal symbols $\{\mathrm{a}, \mathrm{b}, \mathrm{c}, \neg, \land, \lor, (,)\}$ and rules

$$\begin{array}{lll}
\mathrm{F} \to \mathrm{A} & \mathrm{A} \to \mathrm{a} & \mathrm{N} \to \neg\mathrm{F} \\
\mathrm{F} \to \mathrm{N} & \mathrm{A} \to \mathrm{b} & \mathrm{C} \to (\mathrm{F} \land \mathrm{F}) \\
\mathrm{F} \to \mathrm{C} & \mathrm{A} \to \mathrm{c} & \mathrm{D} \to (\mathrm{F} \lor \mathrm{F}) \\
\mathrm{F} \to \mathrm{D} & &
\end{array}$$

Start with F. In each step, replace a left-hand side of a rule with its right-hand side until no more variables are left:

$$\mathrm{F} \Rightarrow \mathrm{N} \Rightarrow \neg\mathrm{F} \Rightarrow \neg\mathrm{D} \Rightarrow \neg(\mathrm{F} \lor \mathrm{F}) \Rightarrow \neg(\mathrm{A} \lor \mathrm{F}) \Rightarrow \neg(\mathrm{b} \lor \mathrm{F}) \Rightarrow \neg(\mathrm{b} \lor \mathrm{A}) \Rightarrow \neg(\mathrm{b} \lor \mathrm{c})$$

C1.2 Alphabets and Formal Languages

Alphabets and Formal Languages

Definition (Alphabets, Words and Formal Languages)

An alphabet Σ is a finite non-empty set of symbols.

A word over Σ is a finite sequence of elements from Σ.
The empty word (the empty sequence of elements) is denoted by ε.
Σ^* denotes the set of all words over Σ.

We write $|w|$ for the length of a word w.

A formal language (over alphabet Σ) is a subset of Σ^*.

German: Alphabet, Zeichen/Symbole, leeres Wort, formale Sprache

Example

$\Sigma = \{\mathtt{a}, \mathtt{b}\}$
$\Sigma^* = \{\varepsilon, \mathtt{a}, \mathtt{b}, \mathtt{aa}, \mathtt{ab}, \mathtt{ba}, \mathtt{bb}, \dots\}$
$|\mathtt{aba}| = 3, |\mathtt{b}| = 1, |\varepsilon| = 0$

Languages: Examples

Example (Languages over $\Sigma = \{\mathtt{a}, \mathtt{b}\}$)

- $S_1 = \{\mathtt{a}, \mathtt{aa}, \mathtt{aaa}, \mathtt{aaaa}, \dots\}$
- $S_2 = \Sigma^*$
- $S_3 = \{\mathtt{a}^n\mathtt{b}^n \mid n \geq 0\} = \{\varepsilon, \mathtt{ab}, \mathtt{aabb}, \mathtt{aaabbb}, \dots\}$
- $S_4 = \{\varepsilon\}$
- $S_5 = \emptyset$
- $S_6 = \{w \in \Sigma^* \mid w$ contains twice as many as as bs$\}$
  $= \{\varepsilon, \mathtt{aab}, \mathtt{aba}, \mathtt{baa}, \dots\}$
- $S_7 = \{w \in \Sigma^* \mid |w| = 3\}$
  $= \{\mathtt{aaa}, \mathtt{aab}, \mathtt{aba}, \mathtt{baa}, \mathtt{bba}, \mathtt{bab}, \mathtt{abb}, \mathtt{bbb}\}$

C1.3 Grammars

Grammars

Definition (Grammars)

A grammar is a 4-tuple $\langle \Sigma, V, P, S \rangle$ with:

1. Σ finite alphabet of terminal symbols
2. V finite set of variables (nonterminal symbols) with $V \cap \Sigma = \emptyset$
3. $P \subseteq (V \cup \Sigma)^+ \times (V \cup \Sigma)^*$ finite set of rules (or productions)
4. $S \in V$ start variable

German: Grammatik, Terminalalphabet, Variablen, Regeln/Produktionen, Startvariable

Rule Sets

What exactly does $P \subseteq (V \cup \Sigma)^+ \times (V \cup \Sigma)^*$ mean?

- $(V \cup \Sigma)^*$: all words over $(V \cup \Sigma)$
- $(V \cup \Sigma)^+$: all non-empty words over $(V \cup \Sigma)$
  in general, for set X: $X^+ = X^* \setminus \{\varepsilon\}$
- $\times$: Cartesian product
- $(V \cup \Sigma)^+ \times (V \cup \Sigma)^*$: set of all pairs $\langle x, y \rangle$, where x non-empty word over $(V \cup \Sigma)$ and y word over $(V \cup \Sigma)$
- Instead of $\langle x, y \rangle$ we usually write rules in the form $x \rightarrow y$.

Rules: Examples

Example

Let $\Sigma = \{\mathtt{a}, \mathtt{b}, \mathtt{c}\}$ and $V = \{\mathsf{X}, \mathsf{Y}, \mathsf{Z}\}$.

Some examples of rules in $(V \cup \Sigma)^+ \times (V \cup \Sigma)^*$:

$$\begin{aligned} \mathsf{X} &\rightarrow \mathsf{X}\mathtt{a}\mathsf{Y} \\ \mathsf{Y}\mathtt{b} &\rightarrow \mathtt{a} \\ \mathsf{XY} &\rightarrow \varepsilon \\ \mathsf{XYZ} &\rightarrow \mathtt{abc} \\ \mathtt{abc} &\rightarrow \mathsf{XYZ} \end{aligned}$$

Derivations

Definition (Derivations)

Let $\langle \Sigma, V, P, S \rangle$ be a grammar. A word $v \in (V \cup \Sigma)^*$ can be derived from word $u \in (V \cup \Sigma)^+$ (written as $u \Rightarrow v$) if

1. $u = xyz$, $v = xy'z$ with $x, z \in (V \cup \Sigma)^*$ and
2. there is a rule $y \rightarrow y' \in P$.

We write: $u \Rightarrow^* v$ if v can be derived from u in finitely many steps (i. e., by using n derivations for $n \in \mathbb{N}_0$).

German: Ableitung

Language Generated by a Grammar

Definition (Languages)

The language generated by a grammar $G = \langle \Sigma, V, P, S \rangle$

$$\mathcal{L}(G) = \{w \in \Sigma^* \mid S \Rightarrow^* w\}$$

is the set of all words from Σ^* that can be derived from S with finitely many rule applications.

German: erzeugte Sprache

Grammars

Examples: blackboard

C1.4 Chomsky Hierarchy

Chomsky Hierarchy

Grammars are organized into the Chomsky hierarchy.

Definition (Chomsky Hierarchy)

- Every grammar is of type 0 (all rules allowed).
- Grammar is of type 1 (context-sensitive)
  if all rules $w_1 \to w_2$ satisfy $|w_1| \leq |w_2|$.
- Grammar is of type 2 (context-free)
  if additionally $w_1 \in V$ (single variable) in all rules $w_1 \to w_2$.
- Grammar is of type 3 (regular)
  if additionally $w_2 \in \Sigma \cup \Sigma V$ in all rules $w_1 \to w_2$.

special case: rule $S \to \varepsilon$ is always allowed if S is the start variable and never occurs on the right-hand side of any rule.

German: Chomsky-Hierarchie, Typ 0, Typ 1 (kontextsensitiv), Typ 2 (kontextfrei), Typ 3 (regulär)

Chomsky Hierarchy

Examples: blackboard

Chomsky Hierarchy

Definition (Type 0–3 Languages)

A language $L \subseteq \Sigma^*$ is of type 0 (type 1, type 2, type 3) if there exists a type-0 (type-1, type-2, type-3) grammar G with $\mathcal{L}(G) = L$.

Type k Language: Example

Example

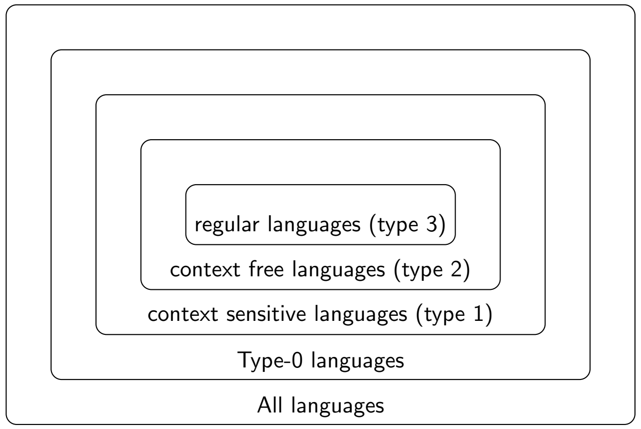
Consider the language L generated by the grammar $\langle \{\mathtt{a}, \mathtt{b}, \mathtt{c}, \neg, \wedge, \vee, (,)\}, \{\mathsf{F}, \mathsf{A}, \mathsf{N}, \mathsf{C}, \mathsf{D}\}, P, \mathsf{F} \rangle$ with the following rules P:

$$\begin{array}{lll} \mathsf{F} \to \mathsf{A} & \mathsf{A} \to \mathtt{a} & \mathsf{N} \to \neg\mathsf{F} \\ \mathsf{F} \to \mathsf{N} & \mathsf{A} \to \mathtt{b} & \mathsf{C} \to (\mathsf{F} \wedge \mathsf{F}) \\ \mathsf{F} \to \mathsf{C} & \mathsf{A} \to \mathtt{c} & \mathsf{D} \to (\mathsf{F} \vee \mathsf{F}) \\ \mathsf{F} \to \mathsf{D} & & \end{array}$$

Questions:

- Is L a type-0 language?
- Is L a type-1 language?
- Is L a type-2 language?
- Is L a type-3 language?

Chomsky Hierarchy


regular languages (type 3)
context free languages (type 2)
context sensitive languages (type 1)
Type-0 languages
All languages

Note: Not all languages can be described by grammars. (Proof?)

C1.5 Summary

Summary

- Languages are sets of symbol sequences.
- Grammars are one possible way to specify languages.
- Language generated by a grammar is the set of all words (of nonterminal symbols) derivable from the start symbol.
- Chomsky hierarchy distinguishes between languages at different levels of expressiveness.

following chapters:

- more about regular languages
- automata as alternative representation of languages