

# Foundations of Artificial Intelligence

## 43. Board Games: Alpha-Beta Search

Thomas Keller

Universität Basel

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May 27, 2016 — 43. Board Games: Alpha-Beta Search

## 43.1 Alpha-Beta Search

### 43.2 Summary

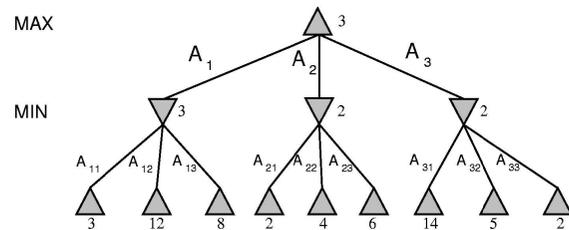
## Board Games: Overview

### chapter overview:

- ▶ 41. Introduction and State of the Art
- ▶ 42. Minimax Search and Evaluation Functions
- ▶ 43. Alpha-Beta Search
- ▶ 44. Monte-Carlo Tree Search: Introduction
- ▶ 45. Monte-Carlo Tree Search: Advanced Topics
- ▶ 46. AlphaGo and Outlook

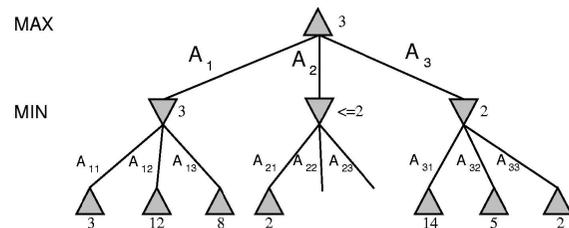
## 43.1 Alpha-Beta Search

## Alpha-Beta Search



Can we save search effort?

We do not need to consider all the nodes!



## Alpha-Beta Search: Generally

Player

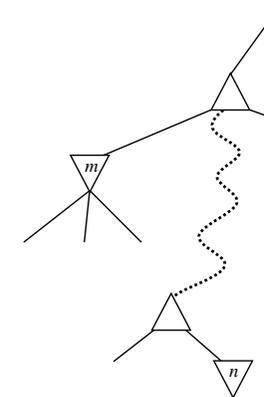
Opponent

..

..

Player

Opponent



If  $m > n$ , then node with utility  $n$  will never be reached when playing perfectly!

## Alpha-Beta Search: Idea

**idea:** Use two values  $\alpha$  and  $\beta$  during minimax depth-first search, such that the following holds for every recursive call:

- ▶ If the utility value in the current subtree is  $\leq \alpha$ , then the subtree **is not interesting** because MAX will never enter it when playing perfectly.
- ▶ If the utility value in the current subtree is  $\geq \beta$ , then the subtree **is not interesting** because MIN will never enter it when playing perfectly.

If  $\alpha \geq \beta$  in the subtree, then the subtree is not interesting and does not have to be searched further ( **$\alpha$ - $\beta$  pruning**).

Starting with  $\alpha = -\infty$  and  $\beta = +\infty$ , alpha-beta search produces the **identical** result as minimax, with lower search effort.

## Alpha-Beta Search: Pseudo Code

- ▶ algorithm skeleton the same as minimax
- ▶ function signature extended by two variables  $\alpha$  and  $\beta$

**function** alpha-beta-main( $p$ )

$\langle v, move \rangle := \text{alpha-beta}(p, -\infty, +\infty)$

**return**  $move$

## Alpha-Beta Search: Pseudo-Code

**function** alpha-beta( $p, \alpha, \beta$ )

**if**  $p$  is final position:

**return**  $\langle u(p), \text{none} \rangle$

initialize  $v$  and  $best\_move$

[as in minimax]

**for each**  $\langle move, p' \rangle \in succ(p)$ :

$\langle v', best\_move' \rangle := \text{alpha-beta}(p', \alpha, \beta)$

  update  $v$  and  $best\_move$

[as in minimax]

**if**  $player(p) = \text{MAX}$ :

**if**  $v \geq \beta$ :

**return**  $\langle v, \text{none} \rangle$

$\alpha := \max\{\alpha, v\}$

**if**  $player(p) = \text{MIN}$ :

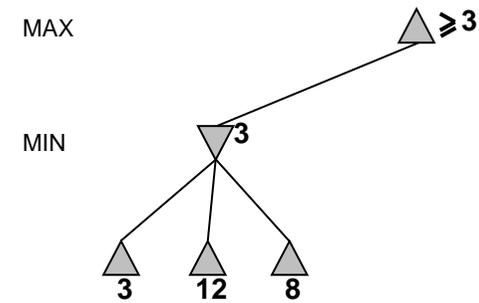
**if**  $v \leq \alpha$ :

**return**  $\langle v, \text{none} \rangle$

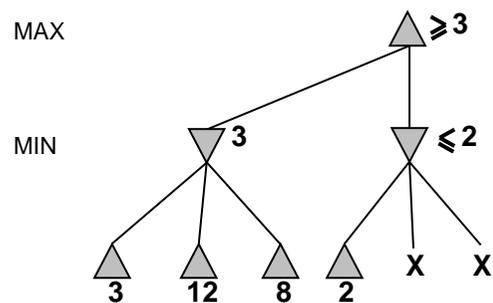
$\beta := \min\{\beta, v\}$

**return**  $\langle v, best\_move \rangle$

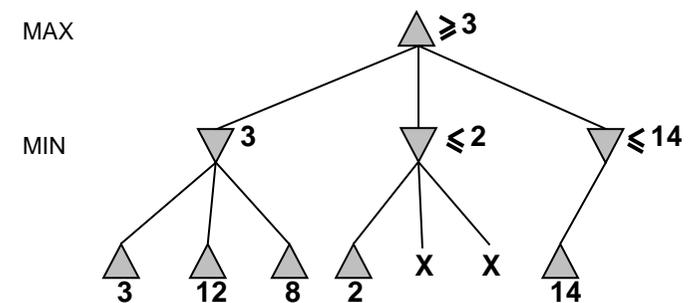
## Alpha-Beta Search: Example



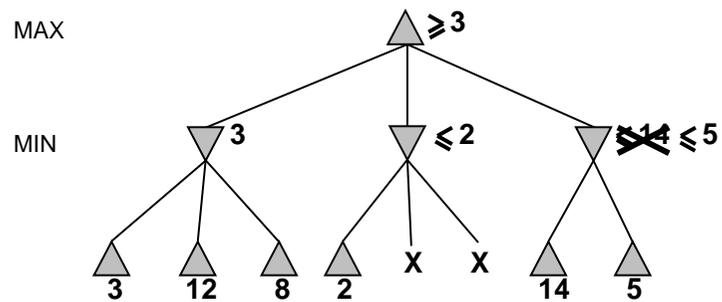
## Alpha-Beta Search: Example



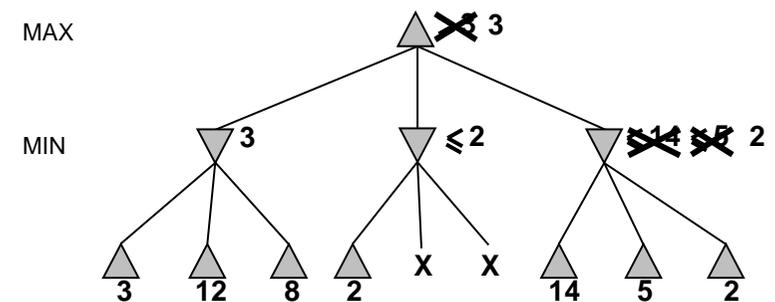
## Alpha-Beta Search: Example



## Alpha-Beta Search: Example



## Alpha-Beta Search: Example



## How Much Do We Get with Alpha-Beta Search?

**assumption:** uniform game tree, depth  $d$ , branching factor  $b \geq 2$ ;  
MAX and MIN positions alternating

- ▶ alpha-beta search most successful if **best move** (for the player whose turn it is) is **considered first**
  - ▶ maximizing move for MAX, minimizing move for MIN

↪ best case: effort reduced from  $O(b^d)$  (minimax) to  $O(b^{d/2})$

↪ doubles the search depth that can be achieved in same time

- ▶ in practice, it is often possible to get close to the optimum

## 43.2 Summary

## Summary

- ▶ **alpha-beta search** stores which utility both players can force somewhere else in game tree and avoids many unnecessary computations
- ▶ in best case, search effort  $O(b^{d/2})$  for uniform trees
- ↪ search depth can be doubled compared to minimax